

# Twists for positroid cells

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# Objects of study: positroid cells

The **matroid** of a  $k \times n$  matrix is the set of subsets of columns which form a basis, written as  $k$ -subsets of  $[n] := \{1, 2, \dots, n\}$ . A **positroid** is the matroid of a 'totally positive matrix': a real-valued matrix whose maximal minors are all non-negative.

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A matroid  $\mathcal{M}$  defines a variety in the  $(k, n)$ -Grassmannian.

$$\Pi(\mathcal{M}) := \{[A] \in Gr(k, n) \mid \text{the matroid of } A \subseteq \mathcal{M}\}$$

## Definition (Positroid cell)

The **positroid cell** of a positroid  $\mathcal{M}$  is

$$\Pi^\circ(\mathcal{M}) := \Pi(\mathcal{M}) - \bigcup_{\text{positroids } \mathcal{M}' \subsetneq \mathcal{M}} \Pi(\mathcal{M}')$$

These cells define a well-behaved stratification of  $Gr(k, n)$ .

# Motivation from double Bruhat cells

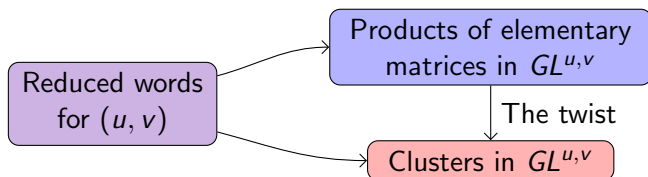
## Example (Double Bruhat cells)

Double Bruhat cells in  $GL(n)$  map to positroid cells under

$$GL(n) \hookrightarrow Gr(n, 2n), \quad A \mapsto [\omega \quad A]$$

where  $\omega$  is the antidiagonal matrix of ones.

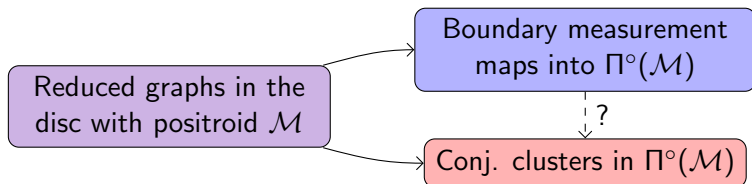
In a double Bruhat cell, the same data indexes two sets of subtori.



Berenstein-Fomin-Zelevinsky introduced a **twist** automorphism of the cell  $GL^{u,v}$  which takes one type of torus to the other.

# The need for a generalized twist

Postnikov described a generalization to all positroid cells.



However, it wasn't proven these constructions defined tori, and the generalized twist was missing for more than a decade!

Then, Marsh-Scott found the twist for the open cell in  $Gr(k, n)$ .

**Our goal!**

Inspired by MS, define the twist automorphism of every  $\Pi^\circ(\mathcal{M})$ .

As a corollary, we prove that the constructions produce tori.

But first, we need to define...

- a **reduced graph** with positroid  $\mathcal{M}$ ,
- its **boundary measurement map**

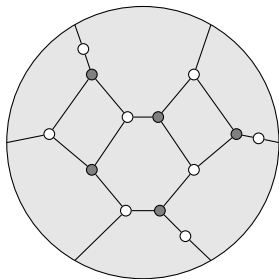
$$\mathbb{B} : \text{a torus} \longrightarrow \Pi^\circ(\mathcal{M}),$$

- and its (conjectural) **cluster**

$$\mathbb{F} : \Pi^\circ(\mathcal{M}) \dashrightarrow \text{a torus}.$$

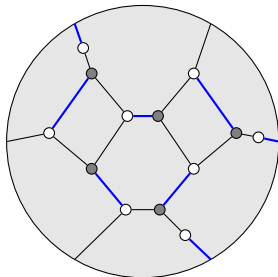
Compared to that, the definition of the twist is elementary.

# Matchings of graphs in the disc



Let  $G$  be a graph in the disc with a 2-coloring of its internal vertices. We assume each boundary vertex is adjacent to one white vertex, and no black or boundary vertices.

A **matching** of  $G$  is a subset of the edges for which every internal vertex is in exactly one edge. We assume a matching exists.



# The positroid of a graph

Observation: Every matching of  $G$  must use

$$k := (\# \text{ of white vertices}) - (\# \text{ of black vertices})$$

boundary vertices.

If we index the boundary vertices clockwise by  $[n]$ , which  $k$ -element subsets of  $[n]$  are the boundary of matchings?

**Theorem (Postnikov, Talaska, Postnikov-Speyer-Williams)**

*The  $k$ -element subsets of  $[n]$  which are the boundary of some matching of  $G$  form a positroid. Every positroid occurs this way.*

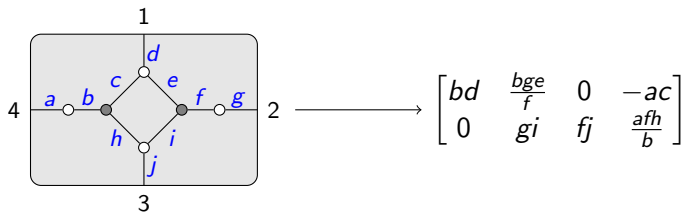
A graph  $G$  is **reduced** if it has the minimal number of faces among all graphs with the same positroid as  $G$ .





# Relations between generating functions

Remarkably, the generating functions satisfy the **Plücker relations**!  
So, for a complex number at each edge of  $G$ , there is a matrix...



...whose  $l$ th maximal minor equals the  $l$ th generating function  $D_l$ ...

$D_{12} = bdgi$	$D_{13} = bdfj$
$D_{14} = adfh$	$D_{23} = begj$
$D_{24} = acgi + aegh$	$D_{34} = acfj$

$\Delta_{12} = bdgi$	$\Delta_{13} = bdfj$
$\Delta_{14} = adfh$	$\Delta_{23} = begj$
$\Delta_{24} = acgi + aegh$	$\Delta_{34} = acfj$

...and this matrix determines a well-defined point in  $Gr(k, n)$ .

# The boundary measurement map

Hence, we have a map

$$\mathbb{C}^{\text{Edges}(G)} \longrightarrow \text{Gr}(k, n)$$

This map is invariant under **gauge transformations**: simultaneously scaling the numbers at each edge incident to a fixed internal vertex.

## Theorem (Postnikov, Talaska, M-Speyer)

For reduced  $G$ , the map  $(\mathbb{C}^*)^{\text{Edges}(G)} \rightarrow \text{Gr}(k, n)$  factors through

$$\mathbb{B} : (\mathbb{C}^*)^{\text{Edges}(E)} / \text{Gauge} \longrightarrow \Pi^\circ(\mathcal{M})$$

where  $\mathcal{M}$  is the positroid of  $G$ .

The map  $\mathbb{B}$  is called the **boundary measurement map**.

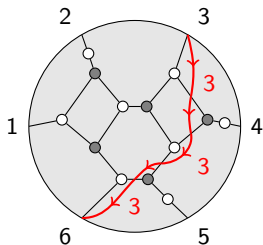
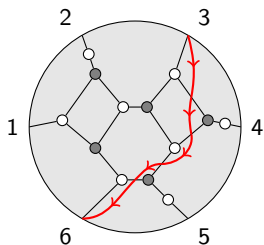
## Conjecture (essentially Postnikov)

The map  $\mathbb{B}$  is an open inclusion.

# Strands in a reduced graph

A **strand** in reduced  $G$  is a path which...

- passes through the midpoints of edges,
- turns right around white vertices,
- turns left around black vertices, and
- begins and ends at boundary vertices.

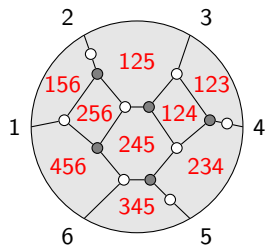


Index a strand by its **source** vertex, and label each face to the **left** of the strand by that label.

# Face labels and the cluster structure

Repeating this for each strand, each face of  $G$  gets **labeled** by a subset of  $[n]$ .

Each face label is a  $k$ -element subset of  $[n]$ , which determines a Plücker coordinate on  $Gr(k, n)$ .



## Conjecture (essentially Postnikov)

*The homogeneous coordinate ring of  $\Pi^\circ(\mathcal{M})$  is a cluster algebra, and the Plücker coordinates of the faces of  $G$  form a cluster.*

The conjecture implies the Plücker coordinates of the faces give a rational map

$$\mathbb{F} : \Pi^\circ(\mathcal{M}) \dashrightarrow (\mathbb{C}^*)^{\text{Faces}(G)} / \text{Scaling}$$

which is an isomorphism on its domain (the **cluster torus**).

# Two conjectural tori associated to a reduced graph

We now see that a **reduced graph**  $G$  with positroid  $\mathcal{M}$  determines two subspaces in  $\Pi^\circ(\mathcal{M})$ , which are both conjecturally tori.

- The image of the **boundary measurement map**

$$\mathbb{B} : (\mathbb{C}^*)^{\text{Edges}(G)} / \text{Gauge} \longrightarrow \Pi^\circ(\mathcal{M})$$

- The domain of definition of the **cluster** of Plücker coordinates

$$\mathbb{F} : \Pi^\circ(\mathcal{M}) \dashrightarrow (\mathbb{C}^*)^{\text{Faces}(G)} / \text{Scaling}$$

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In a simple world, they'd coincide and we'd have an isomorphism

$$\mathbb{F} \circ \mathbb{B} : (\mathbb{C}^*)^{\text{Edges}(G)} / \text{Gauge} \xrightarrow{\sim} (\mathbb{C}^*)^{\text{Faces}(G)} / \text{Scaling}$$

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In the real world, we need a **twist** automorphism  $\tau$  of  $\Pi^\circ(\mathcal{M})$ .

# The twist of a matrix

Let  $A$  be a  $k \times n$  matrix of rank  $k$ , and assume no zero columns. Denote the  $i$ th column of  $A$  by  $A_i$ , with cyclic indices:  $A_{i+n} = A_i$ .

## Definition (The twist)

The **twist**  $\tau(A)$  of  $A$  is the  $k \times n$ -matrix defined on columns by

$$\tau(A)_i \cdot A_i = 1$$

$$\tau(A)_i \cdot A_j = 0, \quad \text{if } A_j \text{ is not in the span of } \{A_i, A_{i+1}, \dots, A_{j-1}\}$$

The columns  $A_i$  and  $\{A_j\}$  in the definition are the 'first' basis of columns encountered when starting at column  $i$  and moving right.

The vector  $\tau(A)_i$  is the left column of the inverse of the submatrix on these columns.

# Example of a twist

## Twisting a matrix

Consider the  $3 \times 4$  matrix

$$A = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 2 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

The first column  $\tau(A)_1$  of the twist is a 3-vector  $v$ , such that...

$$v \cdot \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} = 1, \quad v \cdot \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} = 0, \quad v \cdot \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \text{ is already fixed, and } v \cdot \begin{bmatrix} 0 \\ 2 \\ 1 \end{bmatrix} = 0$$

We see that  $v = \begin{bmatrix} 1 \\ -1 \\ 2 \end{bmatrix}$ . We compute the twist matrix

$$\tau(A) = \begin{bmatrix} 1 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 \\ 2 & 0 & -2 & 1 \end{bmatrix}$$

# The twist on a positroid cell

Twisting matrices descends to a well-defined map of sets

$$Gr(k, n) \xrightarrow{\tau} Gr(k, n)$$

However, this map is not continuous; the defining equations jump when  $A_j$  deforms to a column in the span of  $\{A_i, A_{i+1}, \dots, A_{j-1}\}$ .

## Proposition

*The domains of continuity of  $\tau$  are precisely the positroid cells.*

## Theorem (M-Speyer)

*The twist  $\tau$  restricts to a regular automorphism of  $\Pi^\circ(\mathcal{M})$ .*

The inverse of  $\tau$  is given by a virtually identical formula to  $\tau$ , by reversing the order of the columns.



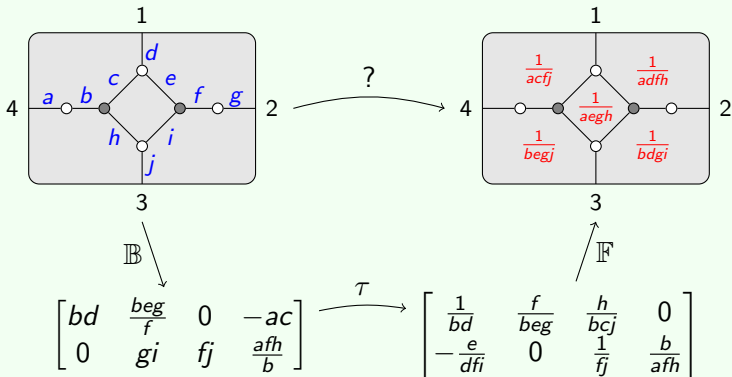
# The induced map on tori

If  $\tau$  takes the image of  $\mathbb{B}$  to the domain of  $\mathbb{F}$ , then the composition

$$\mathbb{F} \circ \tau \circ \mathbb{B} : (\mathbb{C}^*)^{\text{Edges}(G)} / \text{Gauge} \longrightarrow (\mathbb{C}^*)^{\text{Faces}(G)} / \text{Scaling}$$

is a regular morphism. What is this map?

Example: the open cell of  $Gr(2,4)$



# Minimal matchings

It looks like the entries are reciprocals of matchings! Which ones?

## Proposition (Propp)

*The set of matchings of  $G$  with fixed boundary  $I$  has a partial order, with a unique minimal and maximal element (if non-empty).*

## Lemma (M-Speyer)

*If  $I$  is the label of a face in  $G$ , then*

$$\Delta_I \circ \tau \circ \mathbb{B} = \frac{1}{\text{product of edges in } M_I}$$

*where  $M_I$  is the **minimal matching** with boundary  $I$ .*

We have two direct constructions of these minimal matchings.

# The isomorphism of model tori

We can collect these coordinates into a map

$$(\mathbb{C}^*)^{\text{Edges}(G)} \longrightarrow (\mathbb{C}^*)^{\text{Faces}(G)}$$

whose coordinate at a face labeled by  $I$  is the reciprocal of the product of the edges in the minimal matching with boundary  $I$ .

## Lemma (M-Speyer)

*The above map induces an isomorphism of algebraic tori*

$$\mathbb{D} : (\mathbb{C}^*)^{\text{Edges}(E)} / \text{Gauge} \longrightarrow (\mathbb{C}^*)^{\text{Faces}(E)} / \text{Scaling}$$

In fact, the inverse  $\mathbb{D}^{-1}$  can be induced from an explicit map

$$(\mathbb{C}^*)^{\text{Faces}(G)} \longrightarrow (\mathbb{C}^*)^{\text{Edges}(G)}$$

whose coordinate at an edge only uses the two adjacent faces.

# Putting it all together

## Theorem (M-Speyer)

For each reduced graph  $G$ , there is a commutative diagram

$$\begin{array}{ccc}
 (\mathbb{C}^*)^{\text{Edges}(G)} / \text{Gauge} & \begin{array}{c} \xrightarrow{\mathbb{D}} \\ \xleftarrow{\mathbb{D}^{-1}} \end{array} & (\mathbb{C}^*)^{\text{Faces}(G)} / \text{Scaling} \\
 \mathbb{B} \searrow & & \nearrow \mathbb{F} \\
 \Pi^\circ(\mathcal{M}) & \begin{array}{c} \xrightarrow{\tau} \\ \xleftarrow{\tau^{-1}} \end{array} & \Pi^\circ(\mathcal{M})
 \end{array}$$

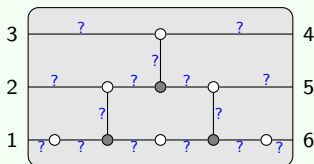
Corollaries:

- The **image** of  $\mathbb{B}$  and the **domain** of  $\mathbb{F}$  are open algebraic tori in  $\Pi^\circ(\mathcal{M})$ , and  $\tau$  takes one to the other.
- The (rational) inverse of  $\mathbb{B}$  is  $\mathbb{D}^{-1} \circ \mathbb{F} \circ \tau$ .
- The (regular) inverse of  $\mathbb{F}$  is  $\tau \circ \mathbb{B} \circ \mathbb{D}^{-1}$ .

# Application: Inverting the boundary measurement map

Let's invert  $\mathbb{B}$  in a classic example!

Example: The unipotent cell in  $GL(3)$ , as a positroid cell



$\mathbb{B}$

$$\begin{bmatrix} 0 & 0 & 1 & a & b & c \\ 0 & -1 & 0 & 0 & d & e \\ 1 & 0 & 0 & 0 & 0 & f \end{bmatrix}$$

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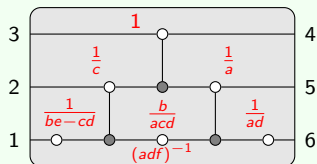
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$$\begin{bmatrix} 0 & 0 & 1 & a & b & c \\ 0 & -1 & 0 & 0 & d & e \\ 1 & 0 & 0 & 0 & 0 & f \end{bmatrix} \xrightarrow{\mathcal{T}} \begin{bmatrix} 0 & 0 & 1 & \frac{1}{a} & \frac{e}{bd-ce} & \frac{1}{c} \\ 0 & -1 & \frac{-b}{d} & \frac{-b}{ad} & \frac{-c}{be-cd} & 0 \\ 1 & \frac{e}{f} & \frac{be-cd}{df} & \frac{be-cd}{adf} & 0 & 0 \end{bmatrix}$$

# Application: Inverting the boundary measurement map

Let's invert  $\mathbb{B}$  in a classic example!

Example: The unipotent cell in  $GL(3)$ , as a positroid cell



$\mathbb{F}$

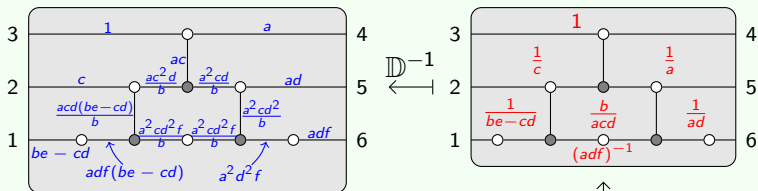
$$\begin{bmatrix} 0 & 0 & 1 & a & b & c \\ 0 & -1 & 0 & 0 & d & e \\ 1 & 0 & 0 & 0 & 0 & f \end{bmatrix} \xrightarrow{\mathcal{T}} \begin{bmatrix} 0 & 0 & 1 & \frac{1}{a} & \frac{e}{bd-ce} & \frac{1}{c} \\ 0 & -1 & \frac{-b}{d} & \frac{-b}{ad} & \frac{-c}{be-cd} & 0 \\ 1 & \frac{e}{f} & \frac{be-cd}{df} & \frac{be-cd}{adf} & 0 & 0 \end{bmatrix}$$



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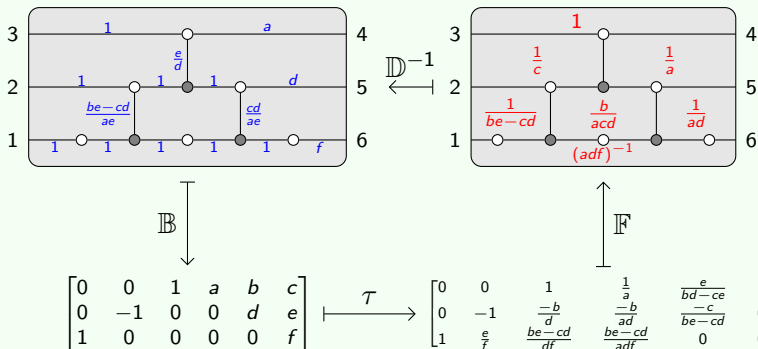


$$\begin{bmatrix} 0 & 0 & 1 & a & b & c \\ 0 & -1 & 0 & 0 & d & e \\ 1 & 0 & 0 & 0 & 0 & f \end{bmatrix} \xrightarrow{\mathcal{T}} \begin{bmatrix} 0 & 0 & 1 & \frac{1}{a} & \frac{e}{bd-ce} & \frac{1}{c} \\ 0 & -1 & \frac{-b}{d} & \frac{-b}{ad} & \frac{-c}{be-cd} & 0 \\ 1 & \frac{e}{f} & \frac{be-cd}{df} & \frac{be-cd}{adf} & 0 & 0 \end{bmatrix}$$

# Application: Inverting the boundary measurement map

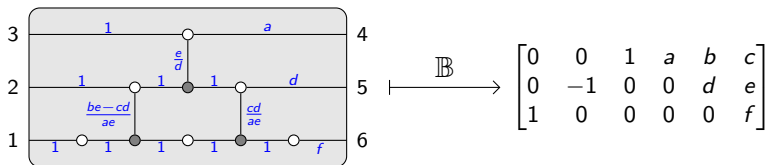
Let's invert  $\mathbb{B}$  in a classic example!

Example: The unipotent cell in  $GL(3)$ , as a positroid cell



# Relation to the Chamber Ansatz

So, we have the following boundary measurement map.



This is equivalent to a factorization into elementary matrices.

$$\begin{bmatrix} a & b & c \\ 0 & d & e \\ 0 & 0 & f \end{bmatrix} = \begin{bmatrix} a & 0 & 0 \\ 0 & d & 0 \\ 0 & 0 & f \end{bmatrix} \begin{bmatrix} 1 & \frac{cd}{ae} & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & \frac{e}{d} \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & \frac{be-cd}{ae} & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Our computation to find this factorization is identical to the *Chamber Ansatz* introduced by Berenstein-Fomin-Zelevinsky.

# Application: Counting matchings

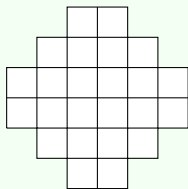
## Corollary

Let  $G$  be a reduced graph with positroid  $\mathcal{M}$ . If  $A$  is a matrix with

- the matroid of  $A$  is contained in  $\mathcal{M}$ , and
- for each face label  $I$ , the minor  $\Delta_I(A) = 1$ ,

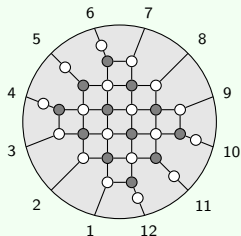
then the maximal minor  $\Delta_J(\tau^{-1}(A))$  counts matchings with boundary  $J$ .

## Example: Domino tilings of the Aztec diamond of order 3



Domino tilings of this shape...

...are the same as matchings of this graph, with boundary  $\{4, 5, 6, 10, 11, 12\}$ .



# Application: Counting matchings

## Example: (continued)

Here is an appropriate  $A$  and its inverse twist.

$$A = \begin{bmatrix} 1 & 6 & 18 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 6 & 1 & 3 & 5 & 0 & -1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 5 & 13 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 6 & 18 & 2 & 2 & 2 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 6 & 2 & 6 & 10 & 0 & -1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 2 & 10 & 26 & 1 & 0 & 0 \end{bmatrix}$$

$$\tau^{-1}(A) = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1 & -2 & -2 & -2 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 6 & 10 & 6 & 2 \\ 0 & 0 & 1 & 0 & 0 & 0 & -1 & -6 & -18 & -26 & -10 & -2 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 5 & 13 & 18 & 6 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & -1 & -3 & -5 & -6 & -1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \end{bmatrix}$$

We compute that  $\Delta_{\{4,5,6,10,11,12\}}(\tau^{-1}(A)) = 64$ . ✓

Finding  $A$  by brute force is probably not efficient, but verifying that a matrix has the necessary properties can be faster than counting.

## Further directions

- Directions to generalize!
  - Positroid cells in  $Gr(k, n) \rightarrow$  projected Richardson cells in partial flag varieties.
  - Reduced graphs in the disc  $\rightarrow$  'reduced graphs' in surfaces.
- If we label strands by their **target** instead of their source, we get a different cluster structure on the same algebra. How are they related?
- Conjecture: The twist is the decategorification of the shift functor in a categorification.  
Any cluster algebra with a Jacobi-finite potential has such a **shift** automorphism. Can this story can be extended?