

Q1)... [12 points] Find a **disjunctive normal form** expression (involving \wedge , \vee , \neg , and P , Q , R) which has the following truth table. Show the steps of your work.

MID I
Solutions

P	Q	R	
T	T	T	T
T	T	F	F
T	F	T	T
T	F	F	F
F	T	T	F
F	T	F	F
F	F	T	T
F	F	F	T

$\leftarrow P \wedge Q \wedge R$
 $\leftarrow P \wedge \neg Q \wedge R$
 $\leftarrow \neg P \wedge \neg Q \wedge R$
 $\leftarrow \neg P \wedge \neg Q \wedge \neg R$

Step ①: Find expressions whose truth tables give a T in the appropriate row & F's elsewhere. There are 4 expressions for this example

Step ②: Take the disjunction of the expressions in step ①.

$$\text{dnf} = (P \wedge Q \wedge R) \vee (P \wedge \neg Q \wedge R) \vee (\neg P \wedge \neg Q \wedge R) \vee (\neg P \wedge \neg Q \wedge \neg R)$$

Find a **conjunctive normal form** expression (involving \wedge , \vee , \neg , and P , Q , R) which has the same truth table above. Show the steps of your work.

Step ①: Negate the output column \longrightarrow

Step ②: Write dnf for this new output col \iff

$$(P \wedge Q \wedge \neg R) \vee (P \wedge \neg Q \wedge \neg R) \vee (\neg P \wedge Q \wedge R) \vee (\neg P \wedge Q \wedge \neg R)$$

F
T
F
T
T
F
F
F

Step ③: Negate this dnf --- use DeMorgan --- get cnf for original table!

$$(\neg P \vee \neg Q \vee R) \wedge (\neg P \vee Q \vee R) \wedge (P \vee \neg Q \vee \neg R) \wedge (P \vee \neg Q \vee R)$$

Q2)... [11 points] Write down the distributive law for \wedge over \vee .

$$A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C)$$

Write down the distributive law for \vee over \wedge .

$$A \vee (B \wedge C) \equiv (A \vee B) \wedge (A \vee C)$$

Write down the two De Morgan laws (involving negations of \wedge and \vee statements).

$$\neg(A \wedge B) \equiv \neg A \vee \neg B$$

$$\neg(A \vee B) \equiv \neg A \wedge \neg B$$

Use the De Morgan and distributive laws to show that the expression

$$[P \wedge (\neg Q) \wedge R] \vee [P \wedge (\neg Q) \wedge (\neg R)] \vee [P \wedge Q \wedge R] \vee \underbrace{\neg[(\neg P) \vee (\neg Q) \vee R]}_{\text{de Morgan}}$$

is logically equivalent to P .

$$\text{Expression} \equiv (P \wedge \neg Q \wedge R) \vee (P \wedge \neg Q \wedge \neg R) \vee (P \wedge Q \wedge R) \vee (P \wedge Q \wedge \neg R)$$

factor out
a "P"
(distributive
Laws)

$$\equiv P \wedge \left[\underbrace{(\neg Q \wedge R) \vee (\neg Q \wedge \neg R)}_{\downarrow \text{factor } \neg Q} \vee \underbrace{(Q \wedge R) \vee (Q \wedge \neg R)}_{\downarrow \text{factor } Q} \right]$$

$$\equiv P \wedge \left[(\neg Q \wedge (R \vee \neg R)) \vee (Q \wedge (R \vee \neg R)) \right]$$

$$\equiv P \wedge \left[(\neg Q \wedge \top) \vee (Q \wedge \top) \right]$$

$$\equiv P \wedge (\neg Q \vee Q)$$

$$\equiv P \wedge \top$$

$$\equiv P \quad \text{done!}$$

Q3]... [12 points] Are the following two expressions logically equivalent. If you say so, please explain why. If you say not, then please give an example which shows that they are different.

$$\forall x [P(x) \rightarrow Q(x)]$$

and

$$(\forall x P(x)) \rightarrow (\forall x Q(x))$$

No

Example (in class!)

Universe = all integers

$P(x)$ = x is even

$Q(x)$ = x is odd

Statement ① is false: eg 2 is even, but not odd.

Statement ② is automatically true, since the hypothesis "every integer is even" is false.

Same question for the expressions

$$\exists x [P(x) \vee Q(x)]$$

and

$$(\exists x P(x)) \vee (\exists x Q(x))$$

Yes

1st \rightarrow 2nd

$\exists x$ such that $P(x) \vee Q(x)$ true
 implies $P(x)$ true or $Q(x)$ true for this value of x
 $\Rightarrow \exists x P(x)$ or $\exists x Q(x)$. $\Rightarrow \exists x P(x) \vee \exists x Q(x)$.

2nd \rightarrow 1st

$$\exists x P(x) \vee \exists x Q(x)$$

$$\Rightarrow \exists x P(x)$$

$$\Rightarrow P(x) \text{ true}$$

$$\Rightarrow P(x) \vee Q(x) \text{ true}$$

$$\rightarrow \exists x (P(x) \vee Q(x))$$

or

$$\exists x Q(x)$$

$$\Rightarrow Q(x) \text{ true}$$

$$\Rightarrow P(x) \vee Q(x) \text{ true}$$

(for a possibly different x than in left column)

Q4]... [15 points] Give a direct proof of the following. If m and n are odd integers, then their product is also odd.

$$m \text{ odd} \Rightarrow m = 2k + 1 \quad \text{for some integer } k.$$

$$n \text{ odd} \Rightarrow n = 2l + 1 \quad \text{for some integer } l.$$

$$\Rightarrow mn = (2k+1)(2l+1) = 4kl + 2k + 2l + 1$$

$$= 2(2kl + k + l) + 1$$

which is of the form $2(\text{integer}) + 1$

\Rightarrow is odd.



Write down the contrapositive of the following statement about integers n . If n^3 is even, then n is also even.

If n is odd, then n^3 is odd.

Prove the statement "If n^3 is even, then n is also even" by giving a proof of its contrapositive.

Start with n is odd

$$\Rightarrow n^2 = n \cdot n = \text{product of 2 odd integers} \\ \text{is odd (by 1st part above)}$$

$$\Rightarrow n^3 = n^2 \cdot n = \text{product of 2 odd integers} \\ \text{is odd (by 1st part above)}$$

$$\Rightarrow n^3 \text{ odd.}$$



Q5]... [15 points] Give a proof of the following: *The cube root of 2 is irrational.* You are free to cite the results of Q4 if they are of any help to you.

Proof by contradiction.

Assume $\sqrt[3]{2}$ is rational.

Thus $\sqrt[3]{2} = \frac{p}{q}$ for $p, q \in \mathbb{Z}^+$. ~~1/4~~

By dividing numerator + denom by some power of 2, we may assume that at least one of p, q is odd. (*)

$$2 = \frac{p^3}{q^3}$$

$$2q^3 = p^3$$

LHS is even $\Rightarrow p^3$ is even

$\Rightarrow p$ is even -- by Q4.

Writing $p = 2k$ for some integer k , we get

$$2q^3 = (2k)^3 = 8k^3$$

$$\Rightarrow q^3 = 4k^3$$

RHS is even $\Rightarrow q^3$ is even

$\Rightarrow q$ is even --- by Q4.

So we have both p and q are even \Rightarrow contradicts (*).

So $\sqrt[3]{2}$ must be irrational. \square

Q6]... [20 points] State the principle of induction.

$P(n)$ = statement involving \oplus integer n .

$$\left. \begin{array}{l} \bullet P(1) \text{ true} \\ \bullet \forall k (P(k) \Rightarrow P(k+1)) \end{array} \right] \Rightarrow P(n) \text{ true } \forall n \in \mathbb{Z}^+$$

Give a proof by induction of the following. For each positive integer n ,

$$P(n) : 1^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$$

Proof

① $P(1)$ is true : $1^2 = 1 \stackrel{??}{=} \frac{1(1+1)(2(1)+1)}{6} = \frac{1(2)(3)}{6} = 1$
 $1 = 1 \checkmark$ true.

② $\forall k [P(k) \Rightarrow P(k+1)]$: Assume $P(k)$ true :

$$1^2 + \dots + k^2 = \frac{k(k+1)(2k+1)}{6}$$

Then

$$\begin{aligned} 1^2 + \dots + k^2 + (k+1)^2 &= \frac{k(k+1)(2k+1)}{6} + (k+1)^2 \\ &= \frac{k(k+1)(2k+1)}{6} + \frac{6(k+1)^2}{6} \\ &= \frac{(k+1)(k(2k+1) + 6(k+1))}{6} \\ &= \frac{(k+1)(2k^2 + 7k + 6)}{6} \\ &= \frac{(k+1)(k+2)(2k+3)}{6} \end{aligned}$$

$$= \frac{(k+1)((k+1)+1)(2(k+1)+1)}{6}$$

& so $P(k+1)$ holds.

By the principle of induction, $P(n)$ true $\forall n \in \mathbb{Z}^+$



Q7)... [15 points] Give a proof by induction of the following. $2^{2n-1} + 3^{2n-1}$ is a multiple of 5 for all integers $n \geq 1$.

$$P(n) : 2^{2n-1} + 3^{2n-1} \text{ is a multiple of } 5$$

$P(1)$ true: $2^{2(1)-1} + 3^{2(1)-1} = 2^1 + 3^1 = 2 + 3 = 5$
is a multiple of 5 ... (5)(1).

$\forall k (P(k) \Rightarrow P(k+1))$: Assume $P(k)$ true.

$$2^{2k-1} + 3^{2k-1} = 5M \text{ for some integer } M.$$

$$\begin{aligned} \text{Now } 2^{2(k+1)-1} + 3^{2(k+1)-1} &= 2^{2k-1+2} + 3^{2k-1+2} \\ &= 2^2 \cdot 2^{2k-1} + 3^2 \cdot 3^{2k-1} \\ &= 4 \cdot 2^{2k-1} + 9 \cdot 3^{2k-1} \\ &\quad \quad \quad \hookrightarrow = (4+5) \\ &= 4(2^{2k-1} + 3^{2k-1}) + 5 \cdot 3^{2k-1} \\ &= 4(5M) + 5 \cdot 3^{2k-1} \quad \dots \text{ by } P(k) \text{ true.} \\ &= 5(4M + 3^{2k-1}) \\ &= \text{multiple of } 5. \Rightarrow P(k+1) \text{ true.} \end{aligned}$$

By induction, $P(n)$ true $\forall n \in \mathbb{Z}^+$.